

Ch 8. Searching and Sorting Arrays

8.1 and 8.3 only

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Definitions of Search and Sort

- Search: find a given item in an array, return the index of the item, or -1 if not found.
- Sort: rearrange the items in an array into some order (smallest to biggest, alphabetical order, etc.).
- There are various methods (algorithms) for carrying out these common tasks.
- Which ones are better? Why?

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Linear Search

- Very simple method.
- Compare first element to target value, if not found then compare second element to target value . . .
- Repeat until:
target value is found (return its index) or
we run out of items (return -1).

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Linear Search in C++

first attempt

```
int searchList (int list[], int size, int target) {  
    int position = -1;           //position of target  
    for (int i=0; i<size; i++)  
    {  
        if (list[i] == target) //found the target!  
            position = i;      //record which item  
    }  
    return position;  
}
```

Is this algorithm correct?

Is this algorithm efficient (does it do unnecessary work)?

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Linear Search in C++

second attempt

```
int searchList (int list[], int size, int value) {
    int index=0;           //index to process the array
    int position = -1;     //position of target
    bool found = false;   //flag, true when target is found

    while (index < size && !found)
    {
        if (list[index] == value) //found the target!
        {
            found = true;         //set the flag
            position = index;     //record which item
        }
        index++;               //increment loop index
    }
    return position;
}
```

Is this algorithm correct?

Is this algorithm efficient (or does it do unnecessary work)?

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Program that uses linear search

```
#include <iostream>
using namespace std;

int searchList(int[], int, int);

int main() {
    const int SIZE=5;
    int idNums[SIZE] = {871, 750, 988, 100, 822};
    int results, id;

    cout << "Enter the employee ID to search for: ";
    cin >> id;

    results = searchList(idNums, SIZE, id);

    if (results == -1) {
        cout << "That id number is not registered\n";
    } else {
        cout << "That id number is found at location ";
        cout << results+1 << endl;
    }
}
```

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Evaluating the Algorithm

- Does it do any unnecessary work?
- Is it efficient? How would we know?
- We measure efficiency of algorithms in terms of number of main steps required to finish.
- For search algorithms, the main step is comparing an array element to the target value.
- Number of steps depends on:
 - size of input array
 - whether or not value is in array
 - where the value is in the array

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Efficiency of Linear Search

how many steps?

N is the number of elements in the array

	N=50,000	In terms of N
Best Case:	1	1
Average Case:	25,000	N/2
Worst Case:	50,000	N

Note: if we search for many items that are not in the array, the average case will be greater than N/2. 8

Binary Search

- Works only for SORTED arrays
- Divide and conquer style algorithm
- Compare target value to middle element in list.
 - if equal, then return its index
 - if less than middle element, repeat the search in the first half of list
 - if greater than middle element, repeat the search in last half of list
- If current search list is narrowed down to 0 elements, return -1

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Binary Search Algorithm

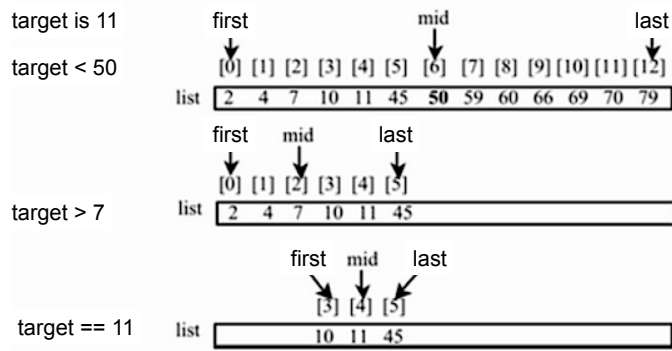
- The algorithm described in pseudocode:


```
while (number of items in list >= 1)
    if (item at middle position is equal to target)
        target is found! End of algorithm
    else
        if (target < middle item)      (narrow search list)
            list = lower half of list
        else
            list = upper half of list
end while

if we reach this point: target not found
```

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Binary Search Algorithm example



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Binary Search in C++

```
int binarySearch (int array[], int size, int target) {
    int first = 0,           //index of (current) first elem
        last = size - 1,    //index of (current) last elem
        middle,             //index of (current) middle elem
        position = -1;      //index of target value
    bool found = false;    //flag

    while (first <= last && !found) {
        middle = (first + last) / 2;    //calculate midpoint

        if (array[middle] == target) {
            found = true;
            position = middle;
        } else if (target < array[middle]) {
            last = middle - 1;          //search lower half
        } else {
            first = middle + 1;         //search upper half
        }
    }
    return position;
}
```

What if first + last is odd?
What if first==last?

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Binary Search

Sample Exam Question!

The target of your search is 42. Given the following array of integers, record the values stored in the variables named `first`, `last`, and `middle` during each iteration of a binary search.

values:	1	7	8	14	20	42	55	67	78	101	112	122	170	179	190
indexes:	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

Repeat the exercise with a target of 82:

first	0	0	4
last	14	6	6
middle	7	3	5

first	0	8	8	8	9
last	14	14	10	8	8
middle	7	11	9	8	X

Note: these are the **indexes**, not the values in the array

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Program using Binary Search

```
#include <iostream>
using namespace std;

int binarySearch(int[], int, int);

int main() {
    const int SIZE=5;
    int idNums[SIZE] = {100, 750, 822, 871, 988};
    int results, id;

    cout << "Enter the employee ID to search for: ";
    cin >> id;

    results = binarySearch(idNums, SIZE, id);

    if (results == -1) {
        cout << "That id number is not registered\n";
    } else {
        cout << "That id number is found at location ";
        cout << results+1 << endl;
    }
}
```

How is this program different from the one on slide 6?

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Efficiency of Binary Search

Calculate worst case for N=1024

Items left to search	Comparisons so far
1024	0
512	1
256	2
128	3
64	4
32	5
16	6
8	7
4	8
2	9
1	10

Goal: calculate this value from N

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$$1024 = 2^{10} \iff \log_2 1024 = 10$$

Efficiency of Binary Search

If N is the number of elements in the array, how many comparisons (steps)?

$$1024 = 2^{10} \iff \log_2 1024 = 10$$

$$N = 2^{\text{steps}} \iff \log_2 N = \text{steps}$$

To what power do I raise 2 to get N?

	N=50,000	In terms of N
Best Case:	1	1
Worst Case:	16	$\log_2 N$

Rounded up to next whole number

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Is $\log_2 N$ better than N ?

Is binary search better than linear search?

Is this really a fair comparison?

Compare values of $N/2$, N , and $\log_2 N$ as N increases:

N	N/2	$\log_2 N$
5	2.5	2.3
50	25	5.6
500	250	9
5,000	2,500	12.3
50,000	25,000	15.6

N and $N/2$ are growing much faster than $\log N$!
 slower growing is more efficient (fewer steps).

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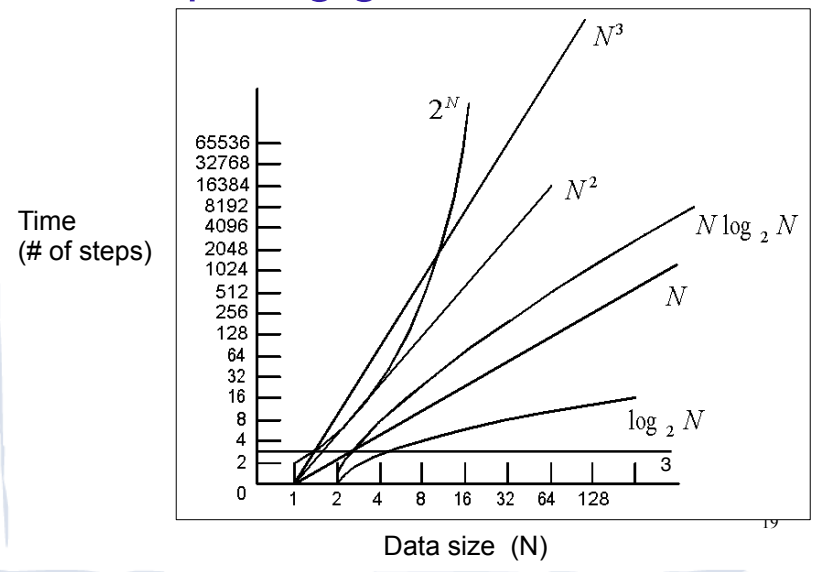
Classifications of (math) functions

Constant	$f(x)=b$	$O(1)$
Logarithmic	$f(x)=\log_b(x)$	$O(\log n)$
Linear	$f(x)=ax+b$	$O(n)$
Linearithmic	$f(x)=x \log_b(x)$	$O(n \log n)$
Quadratic	$f(x)=ax^2+bx+c$	$O(n^2)$
Exponential	$f(x)=b^x$	$O(2^n)$

- Last column is “big Oh notation”, used in CS.
- It ignores all but dominant term, constant factors

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Comparing growth of functions



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Efficiency of Algorithms

- To classify the efficiency of an algorithm:
 - Express “time” (using number of main steps or comparisons), as a function of input size
 - Determine which classification the function fits into.
- Nearer to the top of the classification chart (on slide 18) is slower growth, and more efficient (constant is better than logarithmic, etc.)

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8.3 Sorting Algorithms

- Sort: rearrange the items in an array into ascending or descending order.

- Selection Sort
- Bubble Sort



55 112 78 14 20 179 42 67 190 7 101 1 122 170 8 **unsorted**
1 7 8 14 20 42 55 67 78 101 112 122 170 179 190 **sorted**
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Why is sorting important?

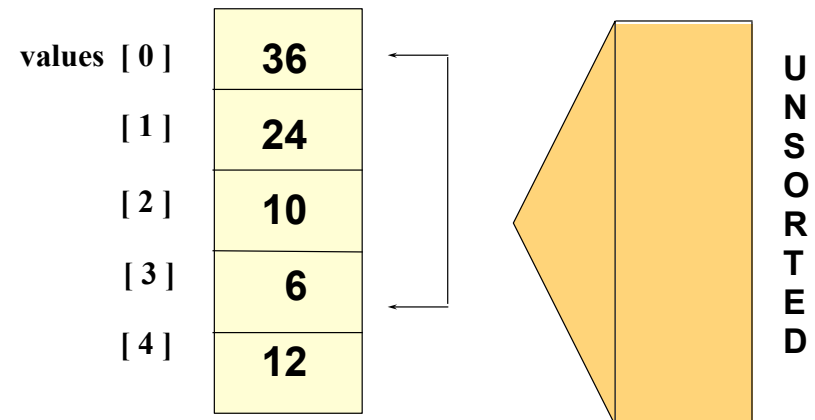
- Searching in a sorted list is much easier than searching in an unsorted list.
- Especially for people:
 - dictionary entries (in a dictionary book)
 - phone book (remember these?)
 - card catalog in library (it used to be drawers of index cards)
 - bank statement: transactions in date order
- Most of the data displayed by computers is sorted.
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Selection Sort

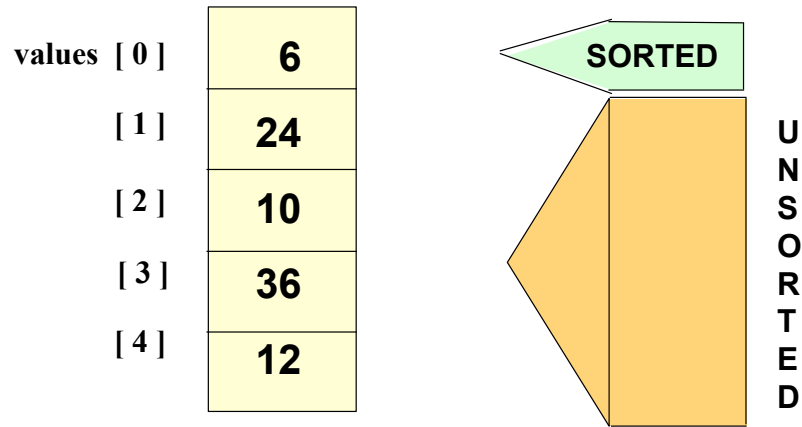
- There is a pass for each position (0..size-1)
- On each pass, the smallest (minimum) element in the rest of the list is exchanged (swapped) with element at the current position.
- The first part of the list (the part that is already processed) is always sorted
- Each pass increases the size of the sorted portion.

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Selection Sort: Pass One

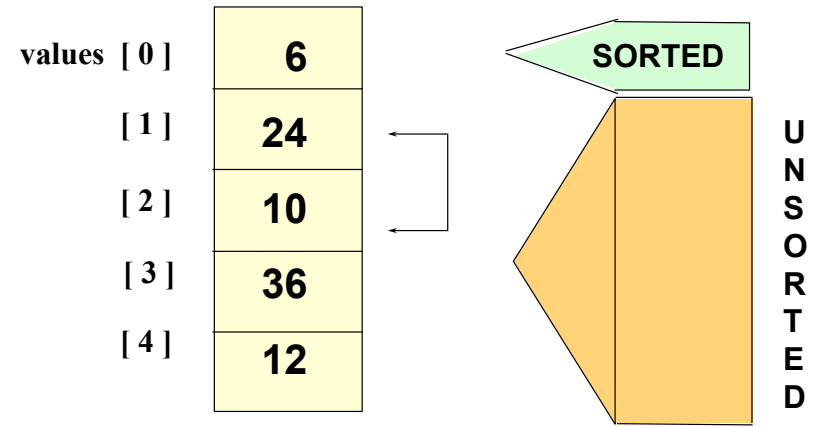


Selection Sort: End Pass One



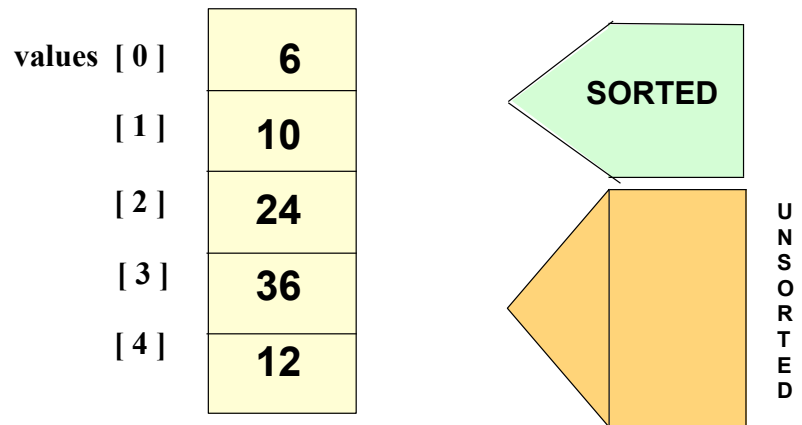
6

Selection Sort: Pass Two



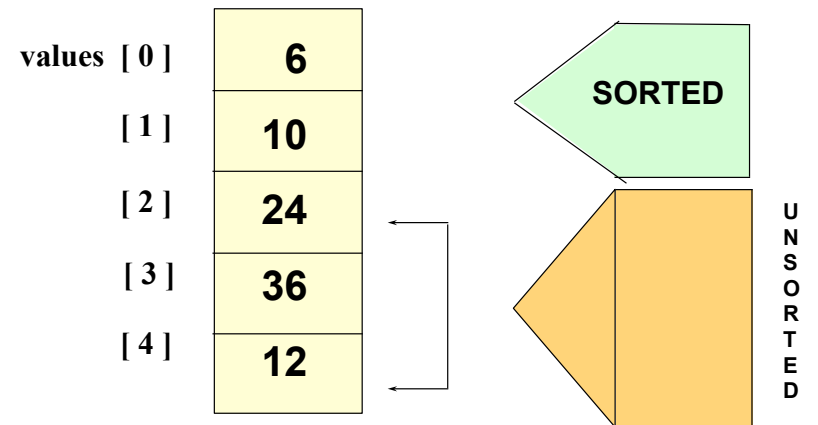
7

Selection Sort: End Pass Two



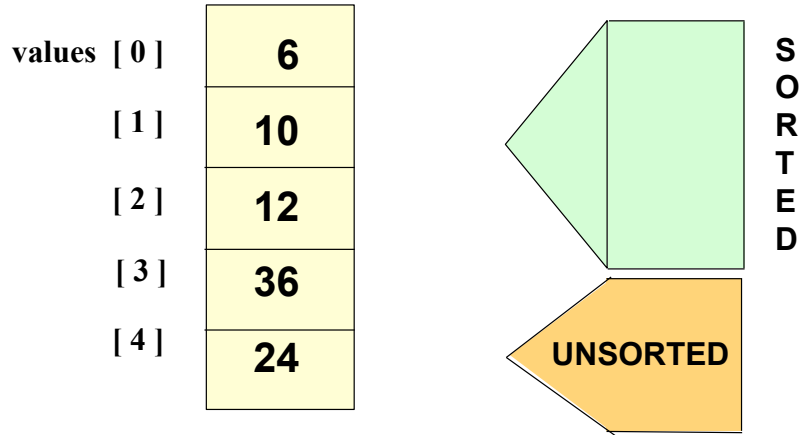
8

Selection Sort: Pass Three



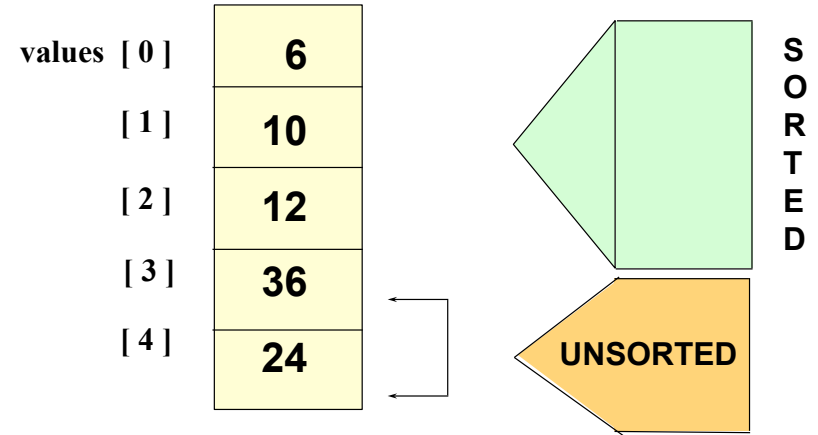
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Selection Sort: End Pass Three



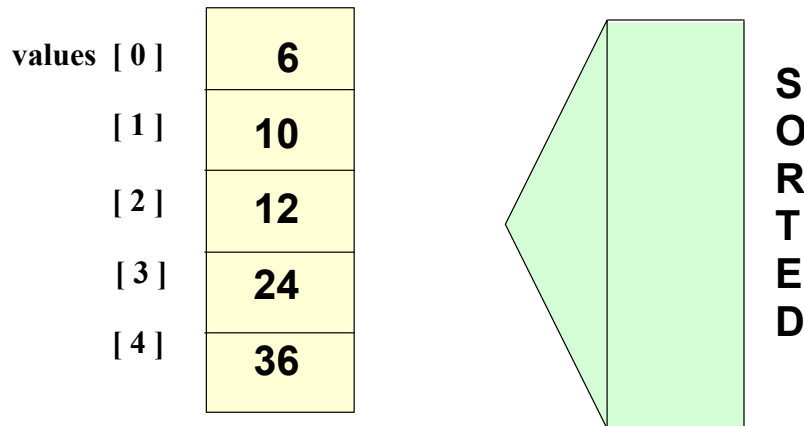
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Selection Sort: Pass Four



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Selection Sort: End Pass Four



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Selection Sort in C++

```
// Returns the index of the smallest element, starting at start
int findIndexOfMin (int array[], int size, int start) {
    int minIndex = start;
    for (int i = start+1; i < size; i++) {
        if (array[i] < array[minIndex]) {
            minIndex = i;
        }
    }
    return minIndex;
}

// Sorts an array, using findIndexOfMin
void selectionSort (int array[], int size) {
    int temp;
    int minIndex;
    for (int index = 0; index < (size -1); index++) {
        minIndex = findIndexOfMin(array, size, index);
        //swap
        temp = array[minIndex];
        array[minIndex] = array[index];
        array[index] = temp;
    }
}
```

Note: saving the index

We need to find the index of the minimum value so that we can do the swap

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Program using Selection Sort

```
#include <iostream>
using namespace std;

int findIndexOfMin (int [], int, int);
void selectionSort(int [], int);
void showArray(int [], int);

int main() {
    int values[6] = {7, 2, 3, 8, 9, 1};

    cout << "The unsorted values are: \n";
    showArray (values, 6);

    selectionSort (values, 6);

    cout << "The sorted values are: \n";
    showArray(values, 6);
}

void showArray (int array[], int size) {
    for (int i=0; i<size; i++)
        cout << array[i] << " ";
    cout << endl;
}
```

Output:

```
The unsorted values are:
7 2 3 8 9 1
The sorted values are:
1 2 3 7 8 9
```

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Efficiency of Selection Sort

- N is the number of elements in the list
- Outer loop (in selectionSort) executes N-1 times
- Inner loop (in minIndex) executes N-1, then N-2, then N-3, ... then once.
- Total number of comparisons (in inner loop):

$$(N-1) + (N-2) + \dots + 2 + 1 = \text{sum of 1 to } N-1$$

$$\text{Note: } N + (N-1) + (N-2) + \dots + 2 + 1 = N(N+1)/2$$

Subtract N from each side:

$$\begin{aligned} (N-1) + (N-2) + \dots + 2 + 1 &= N(N+1)/2 - N \\ &= (N^2+N)/2 - 2N/2 \\ &= (N^2+N-2N)/2 \\ &= N^2/2 - N/2 \end{aligned}$$

O(N²)

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The Bubble Sort

- On each pass:
 - Compare first two elements. If the first is bigger, they exchange places (swap).
 - Compare second and third elements. If second is bigger, exchange them.
 - Repeat until last two elements of the list are compared.
- Repeat this process (keep doing passes) until a pass completes with no exchanges

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Bubble sort

Example: first pass

- 7 2 3 8 9 1 7 > 2, swap
- 2 7 3 8 9 1 7 > 3, swap
- 2 3 7 8 9 1 !(7 > 8), no swap
- 2 3 7 8 9 1 !(8 > 9), no swap
- 2 3 7 8 9 1 9 > 1, swap
- 2 3 7 8 1 9 finished pass 1, did 3 swaps

Note: largest element is now in last position

Note: This is one complete pass!

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Bubble sort

Example: second and third pass

- 2 3 7 **8 1** 9 $2 < 3 < 7 < 8$, no swap, $!(8 < 1)$, swap
- 2 3 7 1 8 9 $(8 < 9)$ no swap
- finished pass 2, did one swap 2 largest elements in last 2 positions

- 2 3 **7 1** 8 9 $2 < 3 < 7$, no swap, $!(7 < 1)$, swap
- 2 3 1 7 8 9 $7 < 8 < 9$, no swap
- finished pass 3, did one swap 3 largest elements in last 3 positions

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Bubble sort

Example: passes 4, 5, and 6

- 2 **3 1** 7 8 9 $2 < 3$, $!(3 < 1)$ swap, $3 < 7 < 8 < 9$
- 2 1 3 7 8 9
- finished pass 4, did one swap
- **2 1** 3 7 8 9 $!(2 < 1)$ swap, $2 < 3 < 7 < 8 < 9$
- 1 2 3 7 8 9
- finished pass 5, did one swap
- 1 2 3 7 8 9 $1 < 2 < 3 < 7 < 8 < 9$, no swaps
- finished pass 6, no swaps, list is sorted!

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Bubble sort

how does it work?

- At the end of the first pass, the largest element is moved to the end (it's bigger than all its neighbors)
- At the end of the second pass, the second largest element is moved to just before the last element.
- The back end (tail) of the list remains sorted.
- Each pass increases the size of the sorted portion.
- No exchanges implies each element is smaller than its next neighbor (so the list is sorted).

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Bubble Sort in C++

```
void bubbleSort (int array[], int size) {  
  
    bool swap;  
    int temp;  
  
    do {  
  
        swap = false;  
        for (int i = 0; i < (size-1); i++) {  
  
            if (array [i] > array[i+1]) {  
  
                temp = array[i];  
                array[i] = array[i+1];  
                array[i+1] = temp;  
                swap = true;  
  
            }  
        }  
    } while (swap);  
}
```

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Program using bubble sort

```
#include <iostream>
using namespace std;

void bubbleSort(int [], int);
void showArray(int [], int);

int main() {
    int values[6] = {7, 2, 3, 8, 9, 1};

    cout << "The unsorted values are: \n";
    showArray (values, 6);

    bubbleSort (values, 6);

    cout << "The sorted values are: \n";
    showArray(values, 6);
}

void showArray (int array[], int size) {
    for (int i=0; i<size; i++)
        cout << array[i] << " ";
    cout << endl;
}
```

Output:

```
The unsorted values are:
7 2 3 8 9 1
The sorted values are:
1 2 3 7 8 9
```

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Efficiency of Bubble Sort

- Each pass makes $N-1$ comparisons
- There will be at most N passes
- So worst case it's: $(N-1)*N = N^2 - N$ **$O(N^2)$**
- If you change the algorithm to look at only the **unsorted** part of the array in each pass, it's exactly like the selection sort:

$$(N-1) + (N-2) + \dots + 2 + 1 = N^2/2 - N/2$$

still $O(N^2)$

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